Docket No. AUS9-2000-0458-US1

LANGUAGE INDEPENDENT MESSAGE MANAGEMENT FOR MULTI-NODE APPLICATION SYSTEMS

BACKGROUND OF THE INVENTION

1. Technical Field:

The present invention relates to an improved networked data processing system and, more particularly, to language translation of messages generated in the system.

2. Description of Related Art:

With the increasing globalization of business, many businesses may find that their computing IT resources are distributed across several countries with servers in one country generating messages and log files in one locale and computers in another country generating messages and log files in a different locale. A locale represents a specific geographical, political or cultural region and encapsulates the information necessary for text to be translated correctly for that locale. Such differences in locale may not be a problem for applications that are executed on a single computer.

However, many software applications are designed such that some processes are executed on one computer while other processes are executed on another computer with the results combined at perhaps still another computer. Each sub-process running on a computer may generate a set of messages and/or log files, perhaps related to errors generated during the execution of the sub-process. Many of these messages and/or log files may

need to be sent back to the originating computer. However, currently, the slave computer executing the sub-processes may not know the locale of the originating computer if the originating computer is more than once removed from the computer executing the sub-processes. Therefore, messages and/or log files generated by the computer executing the sub-processes that are passed back to the originating computer are provided in the locale of the slave computer rather than the locale of the originating computer. Thus, an IT manager may be unable to ascertain key pieces of information due to language barrier.

Therefore, it would be desirable to have a message and log management system for providing information to an originating computer in the locale of the originating computer regardless of the locale of the slave computer.

SUMMARY OF THE INVENTION

The present invention provides a method, system, and computer program product for managing results in a locale independent manner in a multi-node networked data processing system. In one embodiment, a first node sends a command request to a second node. The command request contains a command and a locale in which the text of the result is desired. The first node receives the results of execution from the command request sent to and executed on the second node. The result generated by the second node in response to the command request includes one or more messages, wherein each message contains a unique message identifier, locale in which its associated text is stored, and text associated with the message. Responsive to a determination that the locale of the message text is in a locale different from a desired locale, the first node replaces the message text contained within the result with message text corresponding to the desired locale to produce a modified result and sends the modified result to the requesting client node.

BRIEF DESCRIPTION OF THE DRAWINGS

The novel features believed characteristic of the invention are set forth in the appended claims. The invention itself, however, as well as a preferred mode of use, further objectives and advantages thereof, will best be understood by reference to the following detailed description of an illustrative embodiment when read in conjunction with the accompanying drawings, wherein:

Figure 1 depicts a pictorial representation of a network of data processing systems in which the present invention may be implemented;

Figure 2 depicts a block diagram of a data processing system that may be implemented as a server in accordance with a preferred embodiment of the present invention;

Figure 3 depicts a block diagram illustrating a data processing system in which the present invention may be implemented;

Figure 4 depicts pictorial diagram illustrating a distributed message and logging system in accordance with the present invention;

Figure 5 depicts a block diagram illustrating results arrays for returning results, messages, and log information to a calling node;

Figure 6 depicts a diagram illustrating program flow for message translation in accordance with the present invention; and

Figure 7 depicts a diagram illustrating program flow for manipulating log entries for a distributed application in a multi-node networked data processing system in accordance with the present invention.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENT

With reference now to the figures, **Figure 1** depicts a pictorial representation of a network of data processing systems in which the present invention may be implemented. Network data processing system **100** is a network of computers in which the present invention may be implemented. Network data processing system **100** contains a network **102**, which is the medium used to provide communications links between various devices and computers connected together within network data processing system **100**. Network **102** may include connections, such as wire, wireless communication links, or fiber optic cables.

In the depicted example, a server **104** is connected to network **102** along with storage unit **106**. In addition, clients **108**, **110**, and **112** also are connected to network **102**. These clients **108**, **110**, and **112** may be, for example, personal computers or network computers. In the depicted example, server **104** provides data, such as boot files, operating system images, and applications to clients **108-112**. Clients **108**, **110**, and **112** are clients to server **104**. Network data processing system **100** may include additional servers, clients, and other devices not shown. In the depicted example, network data processing system **100** is the Internet with network **102** representing a worldwide collection of networks and gateways that use the TCP/IP suite of protocols to communicate with one another. At the heart of the Internet is a backbone of high-speed data communication lines between major nodes or host computers, consisting of thousands of commercial,

government, educational and other computer systems that route data and messages. Of course, network data processing system **100** also may be implemented as a number of different types of networks, such as for example, an intranet, a local area network (LAN), or a wide area network (WAN). **Figure 1** is intended as an example, and not as an architectural limitation for the present invention.

Referring to **Figure 2**, a block diagram of a data processing system that may be implemented as a server, such as server **104** in **Figure 1**, is depicted in accordance with a preferred embodiment of the present invention. Data processing system **200** may be a symmetric multiprocessor (SMP) system including a plurality of processors **202** and **204** connected to system bus **206**. Alternatively, a single processor system may be employed. Also connected to system bus **206** is memory controller/cache **208**, which provides an interface to local memory **209**. I/O bus bridge **210** is connected to system bus **206** and provides an interface to I/O bus **212**. Memory controller/cache **208** and I/O bus bridge **210** may be integrated as depicted.

Peripheral component interconnect (PCI) bus bridge **214** connected to I/O bus **212** provides an interface to PCI local bus **216**. A number of modems may be connected to PCI bus **216**. Typical PCI bus implementations will support four PCI expansion slots or add-in connectors. Communications links to network computers **108-112** in **Figure 1** may be provided through modem **218** and network adapter **220** connected to PCI local bus **216** through add-in boards.

Additional PCI bus bridges **222** and **224** provide

interfaces for additional PCI buses **226** and **228**, from which additional modems or network adapters may be supported. In this manner, data processing system **200** allows connections to multiple network computers. A memory-mapped graphics adapter **230** and hard disk **232** may also be connected to I/O bus **212** as depicted, either directly or indirectly.

Those of ordinary skill in the art will appreciate that the hardware depicted in **Figure 2** may vary. For example, other peripheral devices, such as optical disk drives and the like, also may be used in addition to or in place of the hardware depicted. The depicted example is not meant to imply architectural limitations with respect to the present invention.

The data processing system depicted in **Figure 2** may be, for example, an IBM RISC/System 6000 system, a product of International Business Machines Corporation in Armonk, New York, running the Advanced Interactive Executive (AIX) operating system.

With reference now to **Figure 3**, a block diagram illustrating a data processing system is depicted in which the present invention may be implemented. Data processing system **300** is an example of a client computer. Data processing system **300** employs a peripheral component interconnect (PCI) local bus architecture. Although the depicted example employs a PCI bus, other bus architectures such as Accelerated Graphics Port (AGP) and Industry Standard Architecture (ISA) may be used. Processor **302** and main memory **304** are connected to PCI local bus **306** through PCI bridge **308**. PCI bridge **308** also may include an integrated memory controller and cache

memory for processor **302**. Additional connections to PCI local bus **306** may be made through direct component interconnection or through add-in boards. In the depicted example, local area network (LAN) adapter **310**, SCSI host bus adapter **312**, and expansion bus interface **314** are connected to PCI local bus **306** by direct component connection. In contrast, audio adapter **316**, graphics adapter **318**, and audio/video adapter **319** are connected to PCI local bus **306** by add-in boards inserted into expansion slots. Expansion bus interface **314** provides a connection for a keyboard and mouse adapter **320**, modem **322**, and additional memory **324**. Small computer system interface (SCSI) host bus adapter **312** provides a connection for hard disk drive **326**, tape drive **328**, and CD-ROM drive **330**. Typical PCI local bus implementations will support three or four PCI expansion slots or add-in connectors.

An operating system runs on processor **302** and is used to coordinate and provide control of various components within data processing system **300** in **Figure 3**. The operating system may be a commercially available operating system, such as Windows 2000, which is available from Microsoft Corporation. An object oriented programming system such as Java may run in conjunction with the operating system and provide calls to the operating system from Java programs or applications executing on data processing system **300**. "Java" is a trademark of Sun Microsystems, Inc. Instructions for the operating system, the object-oriented operating system, and applications or programs are located on storage devices, such as hard disk drive **326**, and may be loaded into main memory **304** for execution by processor **302**.

Those of ordinary skill in the art will appreciate that the hardware in **Figure 3** may vary depending on the implementation. Other internal hardware or peripheral devices, such as flash ROM (or equivalent nonvolatile memory) or optical disk drives and the like, may be used in addition to or in place of the hardware depicted in **Figure 3**. Also, the processes of the present invention may be applied to a multiprocessor data processing system.

As another example, data processing system **300** may be a stand-alone system configured to be bootable without relying on some type of network communication interface, whether or not data processing system **300** comprises some type of network communication interface. As a further example, data processing system **300** may be a Personal Digital Assistant (PDA) device, which is configured with ROM and/or flash ROM in order to provide non-volatile memory for storing operating system files and/or user-generated data.

The depicted example in **Figure 3** and above-described examples are not meant to imply architectural limitations. For example, data processing system **300** also may be a notebook computer or hand held computer in addition to taking the form of a PDA. Data processing system **300** also may be a kiosk or a Web appliance.

Referring now to **Figures 4-5**, **Figure 4** depicts a pictorial diagram illustrating distributed processing of an application task in accordance with the present invention and **Figure 5** depicts a block diagram illustrating status tables and results fields for managing errors encountered by various nodes in

performing the distributed task in accordance with the present invention. **Figures 4-5** together illustrate the novel distributed data processing result manipulation technique of the present invention. In the present invention, results from a piece of a distributed application executed on one node and returned to an upstream node can be manipulated, for example, by translating messages and/or log files into the locale of a requesting client **402** and by modifying or appending results from one node onto results for the upstream node. Client **402** may be implemented, for example, as data processing system **300** in **Figure 3**. Each of nodes **404-412** may be a server connected to a network, such as network **102** in **Figure 1** and implemented as, for example, data processing system **200** in **Figure 2**.

In a distributed application, client **402** may request a task to be performed by node **404** and in the processing required to complete an application task, node **404** may parcel out sub-tasks for execution on a number of different nodes **406-412**. Errors encountered by one of the parceled processes should be reported on the node **406-412** where the process is executing, on the node **404** from which the application task was initiated, and possibly to the requesting client **402**. Such reporting is enabled for National Language Support (NLS) so that errors can be reported in each node's **402-412** configured locale. The work vended to one node **406-412** is independent of that vended to another **406-412**; all pieces can proceed in parallel and are subject to independent error conditions that may be encountered on each node **406-412**.

A result **502-510** is created on each secondary node **404-412** and sent to the node immediately upstream in the chain of nodes. In the distributed application, the results **504-510** contains result data needed to be used by the upstream node **404** as well as messages, such as, for example, error messages, and also may contain potential log entries.

Since the various nodes **404-412** may use different locales (e.g. U.S. English, U.K. English, Spanish, French, and German) and a client **402** connected to node **404** may use still another locale from some of the nodes **404-412** in the network, messages returned to node **404** are translated into the locale of the client, by node **404**, by retrieving a message text string corresponding to the message ID for the appropriate client locale. If a message does not have a message ID, or if the message ID is unrecognized by node **402**, then node **402** passes the untranslated message text string to the initiating client. Thus, even in the worst case scenario, the initiating client receives some feedback from downstream nodes **404** and **406** even if a translation of the message is not available from node **402**.

As parcels of work (i.e. application sub-tasks) are completed, the results arrays **502-510** are sent back to the originating node. Each results array **502-510** contains entries **511-522**, wherein each entry reflects a result needed by the requesting node or reflects a message or potential log entry. Upon receipt of a results array **502-510**, the receiving node **404-412** may take all, part, or none of the entries from the results array **502-510** and place these entries into an entry in

its own results array **502-510** to be passed to its requesting node. Furthermore, the entries from the received results array **502-510** may be modified, translated, and/or appended to other entries by the receiving node. If an entry contains log information, the receiving node **404-406** may place the log information into its own log file either modified or unmodified and may also pass the log information on to the node upstream from the receiving node as an entry in its own results array.

Thus, in the depicted example, node **406** receives results array (1) **510** from node **410** and results array (2) **508** from node **408**. Results array (1) contains, among other entries, entries **521** and **522** and results array (2) contains, among other entries, entries **519** and **520**. Node **406** takes entry **519** from results array (2) **508** and places it into an entry **518** within its own results array (4) **504**. Node **406** also takes entry **521** from results array (1) **510** and places it within its own results array (4) **504**. Entries **520** and **522** are not placed into results array (4) **504**, however, the information contained within entries **520** and **522** may be utilized by node **406** to generate other messages or results that may be placed into its results array (4) **504**. Also, if entries **520** or **522** contain log information, this information may be added to a log file generated by and stored on node **406** for its own purposes. For example, if the calling node **406** recognizes that one of the log entries is always generated by the secondary node **410** or **408**, the calling node **406** may withhold this information from the calling node **404** rather than writing or passing the information

to the calling node **404**, since this information is unimportant to the calling node **404**.

Similarly, upon receipt of responses **504-506**, node **404** consolidates some, all, or none of the entries **514-518** from results arrays **504** and **506** and adds its own entries to results array **502** which is sent back to the requesting client **402**. In the depicted example, node **404** takes the information from entry **517** and places into its own results array (5) **502** as entry **512**. This information may be entered in entry **512** exactly as it appeared in entry **517** or it may be modified, for example, by appending other information to it or by translating the entry from the locale node **406** into the locale of the requesting client **402**.

Some of the entries **514-518** may contain log data. Node **404** may take this log data and add it to its own log. The log information may also be modified by replacing text generated in one locale with corresponding text matching the locale of the node **404**. Also, the log information may be appended to other log information or a new log entry may be generated based on the information in the log data received from one of the entries **514-518**. To aid in translating entries from one locale to another, each entry **511-522** may include message text and a message ID corresponding to the message or log text. The message ID may be used by the node to identify corresponding translated message text that may be inserted in place of the original message text. These modified log entries may be added to the nodes **404** own log file and/or may be forwarded to the requesting client **402**. However, typically, log entries contain information not useful to

the user of the requesting client **402** and would not be forwarded as an entry in results array (5) **502** to the requesting client **402**.

Similarly, the calling node **404** may write part, all, or none of the log information received from secondary nodes **406-412** into its own log. The calling node **404** may, as desired, write part, none, or all of entries received from nodes **406** and **412** to its results array **502** to be sent to the initiating client **402**. The logging system of the present invention provides a distributed log with log information contained on each node, thus providing each individual node more control over the contents of the log as well as providing a log on the initiating node **404** that may contain only information the node **404** determines is important for the initiating client **402**.

Thus, one aspect of the present invention provides that a single error occurrence is reported to both the node where the error occurs and on the node that originated the distributed work request. On each node, the error is reported to the display terminal and/or in an error log file. The node receiving a result containing log information may write all, part, or none of the log contents to its own log. The node receiving the log information may forward the log information to its calling node in its received form, may modify the log information, or may substitute alternative information within the result. Log entries are recorded as message identifiers plus replacement text plus plain text, allowing log entries to be recorded in the locale of the node where the log is written.

To aid in understanding the present invention, suppose that node **410** is a Japanese node, node **404** is a French node, node **408** if a Spanish node, and node **406** is a German node and consider the following simple example where the nodes **402-408** in the **Figure 4** each generate one message for the result messages, and where the Japanese **410** and French nodes **404** have US English translations installed for their result messages, but where the Spanish **408** and German nodes **406** do not.

The Japanese node **410** creates a result(1) **412** containing message 1000, with replacement parameters 10 and 90, and produces the US English version of the message in the message text field.

The Spanish node **408** creates a result(2) **414** containing message 1211, with replacement parameter 80, and produces the Spanish version of the message in the message text field, since it does not have a US English version available.

The German node **406** creates a result(3) **416** containing message 9415, with replacement parameter C:\, and produces the German version of the message in the message text field, since it does not have a US English version available.

The French node **404** creates a result(4) **418** containing message 1452, with replacement parameter D:\, and produces the US English version of the message in the message text field.

In preparing the message to return to the client **402**, the French node **404** will process each message contained within the result to generate US English messages. In the case of the message generated by the

German node **406**, if the French node **404** does not have a US English translation for the message, then it will pass through the German text. If it does have a translation, US English will be substituted. Likewise for the message for the Spanish node **408**.

Where the nodes **402-410** in an application system share a common message library, the text for a message does not need to be filled in by the node producing the message when it knows there is a node in the return stream that will be able to provide the translation for the message. In the case where the message is not from the common message library, the text should be filled in. This optimization is common for systems that have multiple nodes, and where the nodes have master/slave relationships - the slave can make this optimization with knowledge that its caller has this translation capability.

With reference now to **Figure 6**, a diagram illustrating program flow is depicted in accordance with the present invention. A server, such as server **402** in **Figure 4** receives a request to perform an action from a client node (step **602**). The server sends a request to perform a sub-action needed to carry out the requested action to a second server node within the network, such as, for example, network **100** in **Figure 1** (step **604**). The first server then receives the results, including result messages and/or log messages, from the second server node (step **606**). If the result message(s) and/or log message(s) are in the locale of the client node (step **608**), then the first server does not need to translate the result message(s) and/or log message(s) and merely

sends the result message(s) and/or log message(s) to the client node (step **614**).

If the result message(s) and/or log message(s) are in a locale that is different from the locale of the client node (step **608**), then, using the message ID, the server determines the appropriate translated text that is in the locale of the client node (step **610**). The server then replaces each result message or log message text that is in a locale different from the client node with the appropriate translated text, if that translated text is found (step **612**). The first server then sends the result message(s) and/or log message(s) to the client node with the appropriate result message or log message text replaced with translated text (step **614**). If no replacement text is found, then the first server sends the result message(s) or log message(s) to the client node unmodified.

With reference now to **Figure 7**, a diagram illustrating program flow for manipulating log entries for a distributed application in a multi-node networked data processing system is depicted in accordance with the present invention. To begin, a primary server node in the networked data processing system receives a request to perform an action from a client node (step **702**). The primary server node sends a request to perform a sub-action needed to carry out the requested action to a second server node within the networked data processing system (step **704**). The primary server node then receives the results, including log entry text and ID, from the second server node (step **706**). The primary server node determines which of the log entries should be logged locally and adds the appropriate log entries, if any, to

the local log system (step **708**). Those of ordinary skill in the art will appreciate that the local log system may be implemented in various ways without affecting this invention. Some operating systems include more than one type of log file (within the local log system) based on the type of information to be logged, for example: the system log, the application log and the security log (audit log). The target log to be used is often implied by the log ID (or message identifier). Also note that the log entries added to the local log system may be modified prior to entering them into the local log system. For example, if the log text is in a locale different from the locale of the primary server node, then the log entry text is replaced by appropriate translated log text determined by the log ID. Thus, for example, if the log ID is 20, the primary server node determines the log entry text corresponding to log ID 20 in the locale of the primary server node from a table of log Ids and translation text.

The primary server node then determines which, if any, of the log entries should be sent to the client (step **710**). Any log entries that are to be sent to the client are translated into the locale of the client if the current locale of the log entry is different from the locale of the client (step **712**). This translation is performed using the log ID as described above. The log entries to be sent to the client are placed into the array of results to be returned to the client and then sent to the client (step **714**).

It is important to note that while the present invention has been described in the context of a fully functioning data processing system, those of ordinary

skill in the art will appreciate that the processes of the present invention are capable of being distributed in the form of a computer readable medium of instructions and a variety of forms and that the present invention applies equally regardless of the particular type of signal bearing media actually used to carry out the distribution. Examples of computer readable media include recordable-type media such a floppy disc, a hard disk drive, a RAM, and CD-ROMs and transmission-type media such as digital and analog communications links.

The description of the present invention has been presented for purposes of illustration and description, but is not intended to be exhaustive or limited to the invention in the form disclosed. Many modifications and variations will be apparent to those of ordinary skill in the art. The embodiment was chosen and described in order to best explain the principles of the invention, the practical application, and to enable others of ordinary skill in the art to understand the invention for various embodiments with various modifications as are suited to the particular use contemplated.